

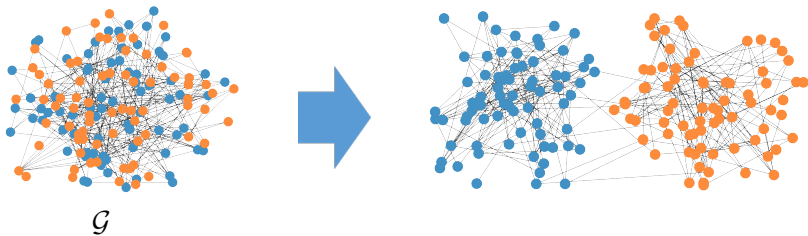
Spectral methods: $\ell_{2,\infty}$ perturbation theory



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Revisit stochastic block model



- Community membership vector

$$x_1^* = \dots = x_{n/2}^* = 1; x_{n/2+1}^* = \dots = x_n^* = -1$$

- observe a graph \mathcal{G} (assuming $p > q$)

$$(i, j) \in \mathcal{G} \text{ with prob. } \begin{cases} p, & \text{if } x_i = x_j \\ q, & \text{else} \end{cases}$$

- **Goal:** recover community memberships $\pm x^*$

Revisit spectral clustering

$$\mathbf{A} = \underbrace{\mathbb{E}[\mathbf{A}]}_{\text{rank 2}} + \mathbf{A} - \mathbb{E}[\mathbf{A}]$$

1. computing the leading eigenvector $\mathbf{u} = [u_i]_{1 \leq i \leq n}$ of $\mathbf{A} - \frac{p+q}{2} \mathbf{1}\mathbf{1}^\top$
2. rounding: output $x_i = \begin{cases} 1, & \text{if } u_i \geq 0 \\ -1, & \text{if } u_i < 0 \end{cases}$

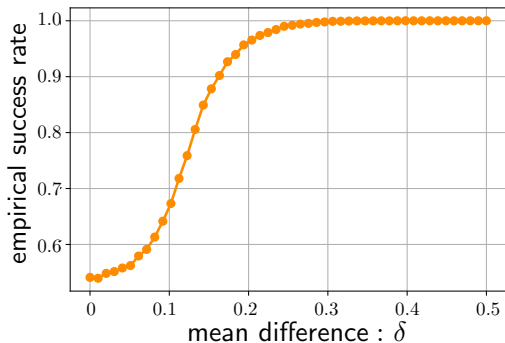
Almost exact recovery

$$\frac{p - q}{\sqrt{p}} \gg \sqrt{\frac{\log n}{n}} \implies \text{almost exact recovery}$$

- Almost exact recovery means

$$\min \left\{ \frac{1}{n} \sum_{i=1}^n \mathbb{1}\{x_i \neq x_i^*\}, \frac{1}{n} \sum_{i=1}^n \mathbb{1}\{x_i \neq -x_i^*\} \right\} = o(1)$$

Empirical performance of spectral clustering



ℓ_2 perturbation theory alone cannot explain exact recovery guarantees

— call for fine-grained analysis

Reverse engineering

Spectral clustering uses signs of u to cluster nodes

Reverse engineering

Spectral clustering uses signs of \mathbf{u} to cluster nodes



It achieves exact recovery iff $u_i u_i^* > 0$ for all i

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A sufficient condition is* $\|\mathbf{u} - \mathbf{u}^*\|_\infty < 1/\sqrt{n}$

Reverse engineering

Spectral clustering uses signs of \mathbf{u} to cluster nodes



It achieves exact recovery iff $u_i u_i^* > 0$ for all i



A sufficient condition is* $\|\mathbf{u} - \mathbf{u}^*\|_\infty < 1/\sqrt{n}$



Need ℓ_∞ perturbation theory

Outline

- An illustrative example: rank-1 matrix denoising
- General ℓ_∞ perturbation theory: symmetric rank-1 case
- Application: exact recovery in community detection
- General $\ell_{2,\infty}$ perturbation theory: rank-r case
- Application: entrywise error in matrix completion

An illustrative example: rank-1 matrix denoising

Setup and algorithm

- Groundtruth: $M^* = \lambda^* \mathbf{u}^* \mathbf{u}^{*\top} \in \mathbb{R}^{n \times n}$, with $\lambda^* > 0$
- Observation: $M = M^* + E$, where E is symmetric, and its upper triangular part comprises of i.i.d. $\mathcal{N}(0, \sigma^2)$ entries
- Estimate \mathbf{u}^* using \mathbf{u} , leading eigenvector of M
- Goal: characterize entrywise error

$$\text{dist}_\infty(\mathbf{u}, \mathbf{u}^*) := \min \{ \|\mathbf{u} - \mathbf{u}^*\|_\infty, \|\mathbf{u} + \mathbf{u}^*\|_\infty \}$$

ℓ_2 guarantees

We start with characterizing noise size

Lemma 6.1

Assume symmetric Gaussian noise model. With high prob., one has

$$\|\mathbf{E}\| \leq 5\sigma\sqrt{n}$$

This in conjunction with Davis-Kahan's $\sin \Theta$ theorem leads to:

$$\text{dist}(\mathbf{u}, \mathbf{u}^*) \leq \frac{2\|\mathbf{E}\|}{\lambda^*} \leq \frac{10\sigma\sqrt{n}}{\lambda^*},$$

as long as $\sigma\sqrt{n} \leq \frac{1-1/\sqrt{2}}{5}\lambda^*$ so that $\|\mathbf{E}\| \leq (1 - 1/\sqrt{2})\lambda^*$

$$\text{--- implies } \text{dist}_\infty(\mathbf{u}, \mathbf{u}^*) \leq \text{dist}(\mathbf{u}, \mathbf{u}^*) \lesssim \frac{\sigma\sqrt{n}}{\lambda}$$

Incoherence

Definition 6.2

Fix a unit vector $\mathbf{u}^* \in \mathbb{R}^n$. Define its incoherence to be

$$\mu := n \|\mathbf{u}^*\|_\infty^2$$

- Range of possible values of μ : $1 \leq \mu \leq n$
- Two extremes: $\mathbf{u}^* = \mathbf{e}_1$, and $\mathbf{u}^* = (1/\sqrt{n}) \cdot \mathbf{1}_n$
- Small μ indicates energy of eigenvector is spread across different entries
- Consider SBM and random Gaussian vectors

ℓ_∞ guarantees for matrix denoising

Theorem 6.3

Suppose that $\sigma\sqrt{n} \leq c_0\lambda^*$ for some sufficiently small constant $c_0 > 0$. Then whp., we have

$$\text{dist}_\infty(\mathbf{u}, \mathbf{u}^*) \lesssim \frac{\sigma(\sqrt{\log n} + \sqrt{\mu})}{\lambda^*}$$

- When $\mu \lesssim \log n$ (i.e., no entries are significantly larger than average), our bound reads

$$\text{dist}_\infty(\mathbf{u}, \mathbf{u}^*) \lesssim \frac{\sigma\sqrt{\log n}}{\lambda^*}$$

- Much tighter than ℓ_2 bound: $\sqrt{n/\log n}$ times smaller

Technical hurdle: dependency

We would like to understand u_l . Since \mathbf{u} is eigenvector of \mathbf{M} , we have

$$\mathbf{M}\mathbf{u} = \lambda\mathbf{u},$$

which yields

$$u_l = \frac{1}{\lambda}[\mathbf{M}]_{l,:}\mathbf{u} = \frac{1}{\lambda}[\mathbf{M}^* + \mathbf{E}]_{l,:}\mathbf{u}$$

\mathbf{u} is dependent on \mathbf{E} ; analyzing $[\mathbf{M}^* + \mathbf{E}]_{l,:}\mathbf{u}$ is challenging

—how to deal with such dependency

An independent proxy

Recall our focus is

$$[\mathbf{M}^* + \mathbf{E}]_{l,:} \mathbf{u}$$

Suppose we have a proxy $\mathbf{u}^{(l)}$ which is **independent** of $[\mathbf{E}]_{l,:}$, then

$$[\mathbf{M}^* + \mathbf{E}]_{l,:} \mathbf{u} = [\mathbf{M}^* + \mathbf{E}]_{l,:} \mathbf{u}^{(l)} + [\mathbf{M}^* + \mathbf{E}]_{l,:} (\mathbf{u} - \mathbf{u}^{(l)})$$

- Independence between $\mathbf{u}^{(l)}$ and $[\mathbf{E}]_{l,:}$
- Proximity between $\mathbf{u}^{(l)}$ and \mathbf{u}

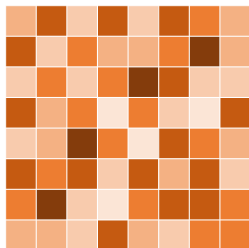
Leave-one-out estimates

For each $1 \leq l \leq n$, construct an auxiliary matrix $M^{(l)}$

$$M^{(l)} := \lambda^* \mathbf{u}^* \mathbf{u}^{*\top} + \mathbf{E}^{(l)},$$

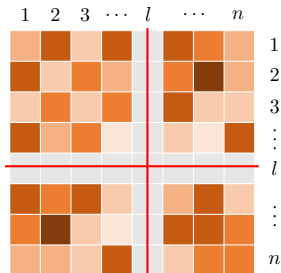
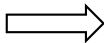
where the noise matrix $\mathbf{E}^{(l)}$ is generated according to

$$E_{i,j}^{(l)} := \begin{cases} E_{i,j}, & \text{if } i \neq l \text{ and } j \neq l, \\ 0, & \text{else.} \end{cases}$$



M

leave one
row/column out



$M^{(l)}$

Leave-one-out estimates (cont.)

For each $1 \leq l \leq n$, construct an auxiliary matrix $\mathbf{M}^{(l)}$

$$\mathbf{M}^{(l)} := \lambda^* \mathbf{u}^* \mathbf{u}^{*\top} + \mathbf{E}^{(l)},$$

where the noise matrix $\mathbf{E}^{(l)}$ is generated according to

$$E_{i,j}^{(l)} := \begin{cases} E_{i,j}, & \text{if } i \neq l \text{ and } j \neq l, \\ 0, & \text{else.} \end{cases}$$

Let $\lambda^{(l)}$ and $\mathbf{u}^{(l)}$ denote respectively leading eigenvalue and leading eigenvector of $\mathbf{M}^{(l)}$

— $\mathbf{u}^{(l)}$ is independent of $[\mathbf{E}]_{l,:}$

Intuition

- Since $\mathbf{u}^{(l)}$ is obtained by dropping only a tiny fraction of data, we expect $\mathbf{u}^{(l)}$ to be extremely close to \mathbf{u} , i.e., $\mathbf{u} \approx \pm \mathbf{u}^{(l)}$
- By construction,

$$\begin{aligned} u_l^{(l)} &= \frac{1}{\lambda^{(l)}} \mathbf{M}_{l,\cdot}^{(l)} \mathbf{u}^{(l)} = \frac{1}{\lambda^{(l)}} \mathbf{M}_{l,\cdot}^* \mathbf{u}^{(l)} = \frac{\lambda^*}{\lambda^{(l)}} u_l^* \mathbf{u}^{*\top} \mathbf{u}^{(l)} \\ &\approx \pm u_l^*. \end{aligned}$$

Proof of Theorem 6.3

What we have learned from ℓ_2 analysis

$$\|\mathbf{E}\| \leq 5\sigma\sqrt{n}$$

$$\text{dist}(\mathbf{u}, \mathbf{u}^*) \leq \frac{10\sigma\sqrt{n}}{\lambda^*}$$

$$|\lambda - \lambda^*| \leq 5\sigma\sqrt{n}$$

$$\max_{j:j \geq 2} |\lambda_j(\mathbf{M})| \leq 5\sigma\sqrt{n}$$

$$\|\mathbf{E}^{(l)}\| \leq \|\mathbf{E}\| \leq 5\sigma\sqrt{n}$$

$$\text{dist}(\mathbf{u}^{(l)}, \mathbf{u}^*) \leq \frac{10\sigma\sqrt{n}}{\lambda^*}$$

$$|\lambda^{(l)} - \lambda^*| \leq 5\sigma\sqrt{n}$$

$$\max_{j:j \geq 2} |\lambda_j(\mathbf{M}^{(l)})| \leq 5\sigma\sqrt{n}$$

Addressing ambiguity

Assume WLOG,

$$\begin{aligned}\|\mathbf{u} - \mathbf{u}^*\|_2 &= \text{dist}(\mathbf{u}, \mathbf{u}^*), \\ \|\mathbf{u}^{(l)} - \mathbf{u}^*\|_2 &= \text{dist}(\mathbf{u}^{(l)}, \mathbf{u}^*), \quad 1 \leq l \leq n\end{aligned}$$

A useful byproduct: if $20\sigma\sqrt{n} < \lambda^*$, then one necessarily has

$$\|\mathbf{u} - \mathbf{u}^{(l)}\|_2 = \text{dist}(\mathbf{u}, \mathbf{u}^{(l)}), \quad 1 \leq l \leq n$$

—*check this*

Bounding $\|u - u^{(l)}\|_2$

Key: view M as perturbation of $M^{(l)}$; apply “sharper” version of Davis-Kahan

$$\|u - u^{(l)}\|_2 \leq \frac{2\|(M - M^{(l)})u^{(l)}\|_2}{\lambda^{(l)} - \max_{j \geq 2} |\lambda_j(M^{(l)})|} \leq \frac{4\|(M - M^{(l)})u^{(l)}\|_2}{\lambda^*}$$

as long as

$$\|M - M^{(l)}\| \leq (1 - 1/\sqrt{2})\left(\lambda^{(l)} - \max_{j \geq 2} |\lambda_j(M^{(l)})|\right),$$

$$\lambda^{(l)} - \max_{j \geq 2} |\lambda_j(M^{(l)})| \geq \lambda^*/2$$

Bounding $\|(M - M^{(l)})\mathbf{u}^{(l)}\|_2$

By design,

$$(M - M^{(l)})\mathbf{u}^{(l)} = \mathbf{e}_l \mathbf{E}_{l,\cdot} \mathbf{u}^{(l)} + u_l^{(l)} (\mathbf{E}_{\cdot,l} - E_{l,l} \mathbf{e}_l),$$

which together with triangle inequality yields

$$\begin{aligned} \|(M - M^{(l)})\mathbf{u}^{(l)}\|_2 &\leq |\mathbf{E}_{l,\cdot} \mathbf{u}^{(l)}| + \|\mathbf{E}_{\cdot,l}\|_2 \cdot |u_l^{(l)}| \\ &\leq 5\sigma \sqrt{\log n} + \|\mathbf{E}_{\cdot,l}\|_2 (|u_l| + \|\mathbf{u} - \mathbf{u}^{(l)}\|_\infty) \\ &\leq 5\sigma \sqrt{\log n} + 5\sigma \sqrt{n} \|\mathbf{u}\|_\infty + 5\sigma \sqrt{n} \|\mathbf{u} - \mathbf{u}^{(l)}\|_2 \end{aligned}$$

Bounding $\|\mathbf{u} - \mathbf{u}^{(l)}\|_2$ (cont.)

Combining previous bounds, we arrive at

$$\begin{aligned}\|\mathbf{u} - \mathbf{u}^{(l)}\|_2 &\leq \frac{20\sigma\sqrt{\log n} + 20\sigma\sqrt{n}\|\mathbf{u}\|_\infty + 20\sigma\sqrt{n}\|\mathbf{u} - \mathbf{u}^{(l)}\|_2}{\lambda^*} \\ &\leq \frac{20\sigma\sqrt{\log n} + 20\sigma\sqrt{n}\|\mathbf{u}\|_\infty}{\lambda^*} + \frac{1}{2}\|\mathbf{u} - \mathbf{u}^{(l)}\|_2,\end{aligned}$$

provided that $40\sigma\sqrt{n} \leq \lambda^*$

Rearranging terms and taking the union bound, we demonstrate that whp.,

$$\|\mathbf{u} - \mathbf{u}^{(l)}\|_2 \leq \frac{40\sigma\sqrt{\log n} + 40\sigma\sqrt{n}\|\mathbf{u}\|_\infty}{\lambda^*} \quad 1 \leq l \leq n$$

Analyzing leave-one-out iterates

Recall that

$$u_i^{(l)} = \frac{1}{\lambda^{(l)}} \mathbf{M}_{i,\cdot}^{(l)} \mathbf{u}^{(l)} = \frac{1}{\lambda^{(l)}} \mathbf{M}_{i,\cdot}^* \mathbf{u}^{(l)} = \frac{\lambda^*}{\lambda^{(l)}} u_i^* \mathbf{u}^{*\top} \mathbf{u}^{(l)}$$

This implies

$$\begin{aligned} u_i^{(l)} - u_i^* &= u_i^* \left(\frac{\lambda^*}{\lambda^{(l)}} \mathbf{u}^{*\top} \mathbf{u}^{(l)} - \mathbf{u}^{*\top} \mathbf{u}^* \right) \\ &= u_i^* \left(\frac{\lambda^* - \lambda^{(l)}}{\lambda^{(l)}} \mathbf{u}^{*\top} \mathbf{u}^{(l)} \right) + u_i^* \mathbf{u}^{*\top} (\mathbf{u}^{(l)} - \mathbf{u}^*) \end{aligned}$$

Analyzing leave-one-out iterates (cont.)

Triangle inequality gives

$$\begin{aligned} |u_l^{(l)} - u_l^*| &\leq |u_l^*| \cdot \frac{|\lambda^* - \lambda^{(l)}|}{|\lambda^{(l)}|} \cdot \|\mathbf{u}^*\|_2 \cdot \|\mathbf{u}^{(l)}\|_2 \\ &\quad + |u_l^*| \cdot \|\mathbf{u}^*\|_2 \cdot \|\mathbf{u}^{(l)} - \mathbf{u}^*\|_2 \\ &\leq |u_l^*| \cdot \frac{10\sigma\sqrt{n}}{\lambda^*} + |u_l^*| \cdot \frac{10\sigma\sqrt{n}}{\lambda^*} \\ &\leq \frac{20\sigma\sqrt{n}}{\lambda^*} \|\mathbf{u}^*\|_\infty \end{aligned}$$

Putting pieces together

Now we come to conclude that

$$\begin{aligned}\|\mathbf{u} - \mathbf{u}^*\|_\infty &= \max_l |u_l - u_l^*| \leq \max_l \left\{ |u_l^{(l)} - u_l^*| + \|\mathbf{u} - \mathbf{u}^{(l)}\|_2 \right\} \\ &\leq \frac{20\sigma\sqrt{n}}{\lambda^*} \|\mathbf{u}^*\|_\infty + \frac{40\sigma\sqrt{\log n} + 40\sigma\sqrt{n}\|\mathbf{u}\|_\infty}{\lambda^*}\end{aligned}$$

One more triangle inequality gives

$$\|\mathbf{u} - \mathbf{u}^*\|_\infty \leq \frac{40\sigma\sqrt{\log n} + 60\sigma\sqrt{n}\|\mathbf{u}^*\|_\infty}{\lambda^*} + \frac{1}{2}\|\mathbf{u} - \mathbf{u}^*\|_\infty,$$

provided that $80\sigma\sqrt{n} \leq \lambda^*$. Rearranging terms yields

$$\|\mathbf{u} - \mathbf{u}^*\|_\infty \leq \frac{80\sigma\sqrt{\log n} + 120\sigma\sqrt{n}\|\mathbf{u}^*\|_\infty}{\lambda^*} = \frac{80\sigma\sqrt{\log n} + 120\sigma\sqrt{\mu}}{\lambda^*},$$

where the last identity results from the definition of μ

General l_∞ perturbation theory

—*symmetric rank-1 case*

Setup and notation

Groundtruth: consider a rank-1 psd matrix $M^* = \lambda^* \mathbf{u}^* \mathbf{u}^{*\top} \in \mathbb{R}^{n \times n}$

Incoherence:

$$\mu := n \|\mathbf{u}^*\|_\infty^2 \quad (1 \leq \mu \leq n)$$

Observations:

$$M = M^* + E \in \mathbb{R}^{n \times n}$$

with E a symmetric noise matrix

Spectral method: return u leading eigenvector of M

Noise assumptions

The entries in the lower triangular part of $\mathbf{E} = [E_{i,j}]_{1 \leq i, j \leq n}$ are independently generated obeying

$$\mathbb{E}[E_{i,j}] = 0, \quad \mathbb{E}[E_{i,j}^2] \leq \sigma^2, \quad |E_{i,j}| \leq B, \quad \text{for all } i \geq j$$

Further, assume that

$$c_b := \frac{B}{\sigma \sqrt{n}/(\mu \log n)} = O(1)$$

Theorem 6.4

With high prob, there exists $z \in \{1, -1\}$ such that

$$\|z\mathbf{u} - \mathbf{u}^*\|_\infty \lesssim \frac{\sigma\sqrt{\mu} + \sigma\sqrt{\log n}}{\lambda^*}, \quad (6.3a)$$

$$\|z\mathbf{u} - \frac{1}{\lambda^*}\mathbf{M}\mathbf{u}^*\|_\infty \lesssim \frac{\sigma\sqrt{\mu}}{\lambda^*} + \frac{\sigma^2\sqrt{n\log n} + \sigma B\sqrt{\mu\log^3 n}}{(\lambda^*)^2} \quad (6.3b)$$

provided that $\sigma\sqrt{n\log n} \leq c_\sigma\lambda^*$ for some sufficiently small constant $c_\sigma > 0$.

- Delocalization of error

First-order expansion

Chain of approximation

$$\mathbf{u} = \frac{M\mathbf{u}}{\lambda} \approx \frac{M\mathbf{u}^*}{\lambda^*} \approx \frac{M^*\mathbf{u}^*}{\lambda^*} = \mathbf{u}^*$$

- first approximation is much tighter than the second one
- important in certain applications such as SBM

**Application: exact recovery in community
detection**

Exact recovery using spectral method

We consider the case when (why?)

$$p = \frac{\alpha \log n}{n}, \quad \text{and} \quad q = \frac{\beta \log n}{n}$$

Theorem 6.5

Fix any constant $\varepsilon > 0$. Suppose $\alpha > \beta > 0$ are sufficiently large, and*

$$(\sqrt{\alpha} - \sqrt{\beta})^2 \geq 2(1 + \varepsilon).$$

With probability $1 - o(1)$, spectral method achieves exact recovery.

Optimality of spectral method

It turns out that when

$$(\sqrt{\alpha} - \sqrt{\beta})^2 \leq 2(1 - \varepsilon),$$

no method whatsoever can achieve exact recovery

—*what's special about $(\sqrt{\alpha} - \sqrt{\beta})^2$ or $(\sqrt{p} - \sqrt{q})^2$?*

Squared Hellinger distance

Definition 6.6

Consider two distributions P and Q over a finite alphabet \mathcal{Y} . The squared Hellinger distance $H^2(P \parallel Q)$ between P and Q is defined as

$$H^2(P \parallel Q) := \frac{1}{2} \sum_{y \in \mathcal{Y}} \left(\sqrt{P(y)} - \sqrt{Q(y)} \right)^2. \quad (6.4)$$

Consider squared Hellinger distance between $\text{Bern}(p)$ and $\text{Bern}(q)$:

$$\begin{aligned} H^2(\text{Bern}(p), \text{Bern}(q)) &:= \frac{1}{2} (\sqrt{p} - \sqrt{q})^2 + \frac{1}{2} (\sqrt{1-p} - \sqrt{1-q})^2 \\ &= (1 + o(1)) \frac{1}{2} (\sqrt{p} - \sqrt{q})^2, \end{aligned}$$

when $p = o(1)$ and $q = o(1)$

Optimality of spectral method (cont.)

The phase transition phenomenon can then be described as

spectral method works if $H^2(\text{Bern}(p), \text{Bern}(q)) \geq (1 + \varepsilon) \frac{\log n}{n}$

no algorithm works if $H^2(\text{Bern}(p), \text{Bern}(q)) \leq (1 - \varepsilon) \frac{\log n}{n}$

for an arbitrary small constant $\varepsilon > 0$

Fine-grained analysis of spectral clustering

Consider “ground-truth” matrix

$$M^* := \mathbb{E}[A] - \frac{p+q}{2} \mathbf{1}\mathbf{1}^\top = \frac{p-q}{2} \begin{bmatrix} \mathbf{1} \\ -\mathbf{1} \end{bmatrix} \begin{bmatrix} \mathbf{1}^\top & -\mathbf{1}^\top \end{bmatrix},$$

which obeys

$$\lambda_1(M^*) := \frac{(p-q)n}{2}, \quad \text{and} \quad \mathbf{u}^* := \frac{1}{\sqrt{n}} \begin{bmatrix} \mathbf{1}_{n/2} \\ -\mathbf{1}_{n/2} \end{bmatrix}.$$

These imply

$$\begin{aligned} \lambda^* &= \frac{n(p-q)}{2}, & \mu &= 1, \\ B &= 1, & \sigma^2 &\leq \max\{p, q\} = p \end{aligned}$$

Invoke ℓ_∞ perturbation theory

ℓ_∞ perturbation bound (6.3b) yields

$$\begin{aligned}\|z\lambda^* \mathbf{u} - \mathbf{M}\mathbf{u}^*\|_\infty &\lesssim \sigma + \frac{\sigma^2 \sqrt{n \log n}}{\lambda^*} + \frac{\sigma B \log^{3/2} n}{\lambda^*} \\ &\leq C \left(\sqrt{p} + \frac{p\sqrt{\log n}}{\sqrt{n}(p-q)} + \frac{\sqrt{p} \log^{3/2} n}{n(p-q)} \right) =: \Delta\end{aligned}$$

for some constant $C > 0$

it boils down to characterizing the entrywise behavior of $\mathbf{M}\mathbf{u}^*$

Bounding entries in $M\mathbf{u}^*$

Lemma 6.7

Suppose that

$$(\sqrt{\alpha} - \sqrt{\beta})^2 \geq 2(1 + \varepsilon)$$

for some quantity $\varepsilon > 0$. Then with probability exceeding $1 - o(1)$, one has

$$\mathbf{M}_{l,\cdot}\mathbf{u}^* \geq \frac{\eta \log n}{\sqrt{n}} \text{ for all } l \leq \frac{n}{2} \text{ and } \mathbf{M}_{l,\cdot}\mathbf{u}^* \leq -\frac{\eta \log n}{\sqrt{n}} \text{ for all } l > \frac{n}{2},$$

where $\eta > 0$ obeys $(\sqrt{\alpha} - \sqrt{\beta})^2 - \eta \log(\alpha/\beta) > 2$.

Key message: entries in $M\mathbf{u}^*$ are bounded away from 0 with correct sign

Completing the picture

On one hand

$$\mathbf{M}_{l,\mathbf{u}^*} \geq \frac{\eta \log n}{\sqrt{n}} \text{ for all } l \leq \frac{n}{2} \text{ and } \mathbf{M}_{l,\mathbf{u}^*} \leq -\frac{\eta \log n}{\sqrt{n}} \text{ for all } l > \frac{n}{2}$$

On the other hand

$$\|z\lambda^*\mathbf{u} - \mathbf{M}\mathbf{u}^*\|_\infty \leq \Delta$$

In sum, if one can show

$$\frac{\eta \log n}{\sqrt{n}} > \Delta \tag{6.5}$$

then it follows that

$$z u_l u_l^* > 0 \text{ for all } 1 \leq l \leq n \implies \text{exact recovery}$$

Proof of relation (6.5)

Our goal is to show

$$\frac{\eta \log n}{\sqrt{n}} \geq C \left(\sqrt{p} + \frac{p\sqrt{\log n}}{\sqrt{n}(p-q)} + \frac{\sqrt{p} \log^{3/2} n}{n(p-q)} \right)$$

- 1st term: $\sqrt{p} \asymp \sqrt{\frac{\log n}{n}} \ll \frac{\eta \log n}{\sqrt{n}}$
- 2nd term: $\frac{p\sqrt{\log n}}{\sqrt{n}(p-q)} \asymp \sqrt{\frac{\log n}{n}} \ll \frac{\eta \log n}{\sqrt{n}}$
- 3rd term: divide discussion into two cases $\alpha/\beta \leq 2$, and $\alpha/\beta \geq 2$

Compare two sets of Bernoullis

Lemma 6.8

Suppose $\alpha > \beta$, $\{W_i\}_{1 \leq i \leq n/2}$ are i.i.d. $\text{Bern}(\frac{\alpha \log n}{n})$, and $\{Z_i\}_{1 \leq i \leq n/2}$ are i.i.d. $\text{Bern}(\frac{\beta \log n}{n})$, which are independent of W_i . For any $t > 0$, one has

$$\mathbb{P} \left(\sum_{i=1}^{n/2} W_i - \sum_{i=1}^{n/2} Z_i \leq t \log n \right) \leq n^{-(\sqrt{a}-\sqrt{b})^2/2+t \log(a/b)/2}.$$

Proof of Lemma 6.7

Note that $M\mathbf{u}^* = (\mathbf{A} - \frac{p+q}{2}\mathbf{1}\mathbf{1}^\top)\mathbf{u}^* = \mathbf{A}\mathbf{u}^*$. Hence

$$M_{1,:}\mathbf{u}^* = \mathbf{A}_{1,:}\mathbf{u}^* = \frac{1}{\sqrt{n}} \left(\sum_{j=1}^{n/2} A_{1,j} - \sum_{j=n/2+1}^n A_{1,j} \right)$$

Apply Lemma 6.8 to obtain with probability at least $1 - n^{-(\sqrt{a}-\sqrt{b})^2/2+\eta\log(a/b)/2} = 1 - o(n^{-1})$

$$M_{1,:}\mathbf{u}^* \geq \frac{\eta \log n}{\sqrt{n}}$$

Invoke union bound to complete proof

Proof of Lemma 6.8

We apply the Laplace transform method: for any $\lambda < 0$

$$\begin{aligned} & \mathbb{P} \left(\sum_{i=1}^{n/2} W_i - \sum_{i=1}^{n/2} Z_i \leq t \log n \right) \\ &= \mathbb{P} \left(\exp \left(\lambda \left(\sum_{i=1}^{n/2} W_i - \sum_{i=1}^{n/2} Z_i \right) \right) \geq \exp(\lambda t \log n) \right) \\ &\leq \frac{\mathbb{E} \left[\exp \left(\lambda \left(\sum_{i=1}^{n/2} W_i - \sum_{i=1}^{n/2} Z_i \right) \right) \right]}{\exp(\lambda t \log n)} \end{aligned}$$

By independence, one has

$$\mathbb{E} \left[\exp \left(\lambda \left(\sum_{i=1}^{n/2} W_i - \sum_{i=1}^{n/2} Z_i \right) \right) \right] = \prod_{i=1}^{n/2} \mathbb{E} [\exp(\lambda W_i)] \mathbb{E} [\exp(-\lambda Z_i)]$$

Proof of Lemma 6.8 (cont.)

By definition and using $1 + x \leq e^x$, one has

$$\begin{aligned}\mathbb{E} [\exp (\lambda W_i)] &= \frac{\alpha \log n}{n} \exp (\lambda) + \left(1 - \frac{\alpha \log n}{n}\right) \\ &\leq \exp \left(\frac{\alpha \log n}{n} \exp (\lambda) - \frac{\alpha \log n}{n}\right)\end{aligned}$$

Similarly for Z_i , one has

$$\mathbb{E} [\exp (-\lambda W_i)] \leq \exp \left(\frac{\beta \log n}{n} \exp (-\lambda) - \frac{\beta \log n}{n}\right)$$

Combine these two to see that

$$\begin{aligned}\mathbb{E} [\exp (\lambda W_i)] \mathbb{E} [\exp (-\lambda Z_i)] \\ \leq \exp \left(\frac{\log n}{n} (\alpha \exp (\lambda) + \beta \exp (-\lambda) - \alpha - \beta)\right)\end{aligned}$$

Proof of Lemma 6.8 (cont.)

Combine previous two pages to see

$$\begin{aligned} \log \mathbb{P} \left(\sum_{i=1}^{n/2} W_i - \sum_{i=1}^{n/2} Z_i \leq t \log n \right) \\ \leq -\lambda t \log n + \frac{n \log n}{2} (\alpha \exp(\lambda) + \beta \exp(-\lambda) - \alpha - \beta) \end{aligned}$$

Set $\lambda = -\log(\alpha/\beta)/2$ to obtain

$$\alpha \exp(\lambda) + \beta \exp(-\lambda) - \alpha - \beta = \alpha \sqrt{\frac{\beta}{\alpha}} + \beta \sqrt{\frac{\alpha}{\beta}} - \alpha - \beta = -(\sqrt{\alpha} - \sqrt{\beta})^2$$

and proof is finished

General $l_{2,\infty}$ perturbation theory

—rank- r case

Setup and notation

Groundtruth: consider a rank- r matrix $M^* = U^* \Sigma^* V^{*\top} \in \mathbb{R}^{n_1 \times n_2}$, with singular values $\sigma_1^* \geq \sigma_2^* \geq \dots \geq \sigma_r^* > 0$ (assume $n_1 \leq n_2$)

Two convenient notation:

$$\kappa := \frac{\sigma_1^*}{\sigma_r^*}, \quad n := n_1 + n_2$$

Observations:

$$M = M^* + E \in \mathbb{R}^{n_1 \times n_2}$$

with E a noise matrix

Spectral method: return U, V where $M = U \Sigma V^\top + U_\perp \Sigma_\perp V_\perp^\top$

Noise assumptions

The entries in $\mathbf{E} = [E_{i,j}]_{1 \leq i \leq n_1, 1 \leq j \leq n_2}$ are independently generated obeying

$$\mathbb{E}[E_{i,j}] = 0, \quad \mathbb{E}[E_{i,j}^2] \leq \sigma^2, \quad |E_{i,j}| \leq B, \quad \text{for all } i, j$$

Further, assume that

$$c_b := \frac{B}{\sigma \sqrt{n_1 / (\mu \log n)}} = O(1)$$

$\ell_{2,\infty}$ distance between U and U^*

Need to take into account rotation ambiguity

—*which rotation matrix to use?*

Definition 6.9

For any square matrix Z with SVD $Z = U_Z \Sigma_Z V_Z^T$, define

$$\text{sgn}(Z) := U_Z V_Z^T \quad (6.6)$$

to be the matrix sign function of Z .

Use $\text{sgn}(U^T U^*)$ —solution to procrustes problem, which yields

$$\|U \text{sgn}(U^T U^*) - U^*\|_{2,\infty}$$

Incoherence of subspace

Definition 6.10

Fix an orthonormal matrix $U^* \in \mathbb{R}^{n \times r}$. Define its incoherence to be

$$\mu(U^*) := \frac{n \|U^*\|_{2,\infty}^2}{r}$$

—recover incoherence of eigenvector when $r = 1$

- For $M^* = U^* \Sigma^* V^{*\top}$, define $\mu(M^*) := \max\{\mu(U^*), \mu(V^*)\}$

$\ell_{2,\infty}$ perturbation theory

Define $\mathbf{H}_U := \mathbf{U}^\top \mathbf{U}^*$ and $\mathbf{H}_V := \mathbf{V}^\top \mathbf{V}^*$

Theorem 6.11

With probability at least $1 - O(n^{-5})$, one has

$$\begin{aligned} \max \left\{ \|\mathbf{U} \operatorname{sgn}(\mathbf{H}_U) - \mathbf{U}^*\|_{2,\infty}, \|\mathbf{V} \operatorname{sgn}(\mathbf{H}_V) - \mathbf{V}^*\|_{2,\infty} \right\} \\ \lesssim \frac{\sigma \sqrt{r} (\kappa \sqrt{\frac{n_2}{n_1}} \mu + \sqrt{\log n})}{\sigma_r^*}, \end{aligned}$$

provided that $\sigma \sqrt{n \log n} \leq c_1 \sigma_r^*$ for some sufficiently small constant $c_1 > 0$.

Entrywise reconstruction error

Recall $M = U\Sigma V^\top + U_\perp \Sigma_\perp V_\perp^\top$

Corollary 6.12

In addition, if $\sigma\kappa\sqrt{n\log n} \leq c_2\sigma_r^$ for some small enough constant $c_2 > 0$, then the following holds with probability at least $1 - O(n^{-5})$:*

$$\|U\Sigma V^\top - M^*\|_\infty \lesssim \sigma\kappa^2\mu r \sqrt{\frac{(n_2/n_1)\log n}{n_1}}$$

De-localization of estimation error

For simplicity, let us consider the case where $\mu, \kappa, n_2/n_1 = O(1)$. Davis-Kahan theorem results in the following ℓ_2 estimation guarantees

$$\text{dist}_{\mathbb{F}}(\mathbf{U}, \mathbf{U}^*) \leq \sqrt{r} \text{dist}(\mathbf{U}, \mathbf{U}^*) \lesssim \frac{\sigma \sqrt{nr}}{\sigma_r^*}$$

In comparison, the $\ell_{2,\infty}$ bound derived in Theorem 6.11 simplifies to

$$\min_{\mathbf{R} \in \mathcal{O}^{r \times r}} \|\mathbf{U}\mathbf{R} - \mathbf{U}^*\|_{2,\infty} \leq \|\mathbf{U} \text{sgn}(\mathbf{H}) - \mathbf{U}^*\|_{2,\infty} \lesssim \frac{\sigma \sqrt{r \log n}}{\sigma_r^*}$$

De-localization of estimation error (cont.)

For the matrix reconstruction error, one has

$$\|U\Sigma V^T - M^*\| \leq 2\|M - M^*\| \lesssim \sigma\sqrt{n},$$

which implies $\|U\Sigma V^T - M^*\|_F \lesssim \sigma\sqrt{nr}$

In comparison, one has

$$\|U\Sigma V^T - M^*\|_\infty \lesssim \sigma r \sqrt{\frac{\log n}{n}}$$

Application: entrywise error in matrix completion

Low-rank matrix completion

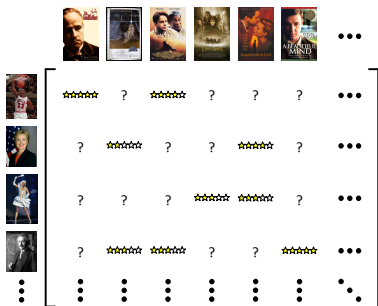


figure credit: Candès

- consider a low-rank matrix $M^* = U^* \Sigma^* V^{*\top}$
- each entry $M_{i,j}^*$ is observed independently with prob. p
- **intermediate goal:** estimate U^*, V^*

Spectral method for matrix completion

1. identify the key matrix M^*
2. construct surrogate matrix $M \in \mathbb{R}^{n \times n}$ as

$$M_{i,j} = \begin{cases} \frac{1}{p} M_{i,j}^*, & \text{if } M_{i,j}^* \text{ is observed} \\ 0, & \text{else} \end{cases}$$

- **rationale for rescaling:** ensures $\mathbb{E}[M] = M^*$

3. compute the rank- r SVD $U\Sigma V^\top$ of M , and return (U, Σ, V)

ℓ_2 guarantees for matrix completion

Theorem 6.13

Suppose that $n_1 p \geq C_1 \kappa^2 \mu r \log n_2$ for some sufficiently large constant $C_1 > 0$. Then with probability exceeding $1 - O(n_2^{-10})$,

$$\max \left\{ \text{dist}(\mathbf{U}, \mathbf{U}^*), \text{dist}(\mathbf{V}, \mathbf{V}^*) \right\} \lesssim \kappa \sqrt{\frac{\mu r \log n_2}{n_1 p}}.$$

- Key: bound $\|\mathbf{M} - \mathbf{M}^*\|$ by $\sqrt{\frac{\mu r \log n_2}{n_1 p}} \|\mathbf{M}^*\|$ (homework)

$\ell_{2,\infty}$ guarantees for matrix completion

Theorem 6.14

Suppose that $n_1 \leq n_2$ and $n_1 p \geq C \kappa^4 \mu^2 r^2 \log n$ for some sufficiently large constant $C > 0$. Then with high prob., we have

$$\begin{aligned} \max\{\|U \operatorname{sgn}(H_U) - U^*\|_{2,\infty}, \|V \operatorname{sgn}(H_V) - V^*\|_{2,\infty}\} \\ \leq \kappa^2 \sqrt{\frac{\mu^3 r^3 \log n}{n_1^2 p}}; \\ \|U \Sigma V^\top - M^*\|_\infty \lesssim \kappa^2 \mu^2 r^2 \sqrt{\frac{\log n}{n_1^3 p}} \|M^*\| \end{aligned}$$

Proof of Theorem 6.14

Recall our notation $\mathbf{E} = \mathbf{M} - \mathbf{M}^* = p^{-1}\mathcal{P}_\Omega(\mathbf{M}^*) - \mathbf{M}^*$. It is straightforward to check that \mathbf{E} satisfies noise assumptions with

$$\sigma^2 := \frac{\|\mathbf{M}^*\|_\infty^2}{p}, \quad \text{and} \quad B := \frac{\|\mathbf{M}^*\|_\infty}{p}$$

In addition, from the relation $B = c_b \sigma \sqrt{n_1 / (\mu \log n)}$, it is seen that $c_b = O(1)$ holds as long as $n_1 p \gtrsim \mu \log n$.

With these preparations in place, the claims in Theorem 6.14 follow directly from Theorem 6.11 and

$$\|\mathbf{M}^*\|_\infty \leq \mu r \|\mathbf{M}^*\| / \sqrt{n_1 n_2}$$

What we have not discussed so far

- More applications of spectral methods
- Uncertainty quantification for spectral estimators
- Precise asymptotic analysis of spectral estimators
- Variants of spectral methods with certain advantages
- ℓ_p analysis of spectral methods